

7.5 The Conjugate Gradient Method

The conjugate gradient method of Hestenes and Stiefel [HS] was originally developed as a direct method designed to solve an $n \times n$ positive definite linear system. As a direct method it is generally inferior to Gaussian elimination with pivoting since both methods require n steps to determine a solution, and the steps of the conjugate gradient method are more computationally expensive than those in Gaussian elimination.

However, the conjugate gradient method is very useful when employed as an iterative approximation method for solving large sparse systems with nonzero entries occurring in predictable patterns. These problems frequently arise in the solution of boundary-value problems. When the matrix has been preconditioned to make the calculations more effective, good results are obtained in only about \sqrt{n} steps. Employed in this way, the method is preferred over Gaussian elimination and the previously-discussed iterative methods.

Throughout this section we assume that the matrix A is positive definite. We will use the *inner product* notation

$$\langle \mathbf{x}, \mathbf{y} \rangle = \mathbf{x}^t \mathbf{y}, \tag{7.25}$$

where \mathbf{x} and \mathbf{y} are n -dimensional vectors. We will also need some additional standard results from linear algebra. A review of this material is found in Section 9.1.

The next result follows easily from the properties of transposes (see Exercise 12).

Theorem 7.30 For any vectors \mathbf{x} , \mathbf{y} , and \mathbf{z} and any real number α , we have

- (i) $\langle \mathbf{x}, \mathbf{y} \rangle = \langle \mathbf{y}, \mathbf{x} \rangle$;
- (ii) $\langle \alpha \mathbf{x}, \mathbf{y} \rangle = \langle \mathbf{x}, \alpha \mathbf{y} \rangle = \alpha \langle \mathbf{x}, \mathbf{y} \rangle$;
- (iii) $\langle \mathbf{x} + \mathbf{z}, \mathbf{y} \rangle = \langle \mathbf{x}, \mathbf{y} \rangle + \langle \mathbf{z}, \mathbf{y} \rangle$;
- (iv) $\langle \mathbf{x}, \mathbf{x} \rangle \geq 0$;
- (v) $\langle \mathbf{x}, \mathbf{x} \rangle = 0$ if and only if $\mathbf{x} = \mathbf{0}$. ■

When A is positive definite, $\langle \mathbf{x}, A\mathbf{x} \rangle = \mathbf{x}^t A \mathbf{x} > 0$ unless $\mathbf{x} = \mathbf{0}$. Also, since A is symmetric, we have $\mathbf{x}^t A \mathbf{y} = \mathbf{x}^t A^t \mathbf{y} = (A\mathbf{x})^t \mathbf{y}$, so in addition to the results in Theorem 7.30, we have for each \mathbf{x} and \mathbf{y} ,

$$\langle \mathbf{x}, A\mathbf{y} \rangle = \langle A\mathbf{x}, \mathbf{y} \rangle. \tag{7.26}$$

The following result is a basic tool in the development of the conjugate gradient method.

Theorem 7.31 The vector \mathbf{x}^* is a solution to the positive definite linear system $A\mathbf{x} = \mathbf{b}$ if and only if \mathbf{x}^* minimizes

$$g(\mathbf{x}) = \langle \mathbf{x}, A\mathbf{x} \rangle - 2\langle \mathbf{x}, \mathbf{b} \rangle. \quad \blacksquare$$

or the coefficients of

630
-12600
56700
-88200
44100

r to that estimated in

ve the estimate using

rices:

$\times 3$ Hilbert matrix H .

Proof Let \mathbf{x} and $\mathbf{v} \neq \mathbf{0}$ be fixed vectors and t a real number variable. We have

$$\begin{aligned} g(\mathbf{x} + t\mathbf{v}) &= \langle \mathbf{x} + t\mathbf{v}, A\mathbf{x} + tA\mathbf{v} \rangle - 2\langle \mathbf{x} + t\mathbf{v}, \mathbf{b} \rangle \\ &= \langle \mathbf{x}, A\mathbf{x} \rangle + t\langle \mathbf{v}, A\mathbf{x} \rangle + t\langle \mathbf{x}, A\mathbf{v} \rangle + t^2\langle \mathbf{v}, A\mathbf{v} \rangle - 2\langle \mathbf{x}, \mathbf{b} \rangle - 2t\langle \mathbf{v}, \mathbf{b} \rangle \\ &= \langle \mathbf{x}, A\mathbf{x} \rangle - 2\langle \mathbf{x}, \mathbf{b} \rangle + 2t\langle \mathbf{v}, A\mathbf{x} \rangle - 2t\langle \mathbf{v}, \mathbf{b} \rangle + t^2\langle \mathbf{v}, A\mathbf{v} \rangle, \end{aligned}$$

so

$$g(\mathbf{x} + t\mathbf{v}) = g(\mathbf{x}) + 2t\langle \mathbf{v}, A\mathbf{x} - \mathbf{b} \rangle + t^2\langle \mathbf{v}, A\mathbf{v} \rangle. \quad (7.27)$$

Since \mathbf{x} and \mathbf{v} are fixed, we can define the quadratic function h in t by

$$h(t) = g(\mathbf{x} + t\mathbf{v}).$$

Then h assumes a minimal value when $h'(t) = 0$, because its t^2 coefficient, $\langle \mathbf{v}, A\mathbf{v} \rangle$, is positive. Since

$$h'(t) = 2\langle \mathbf{v}, A\mathbf{x} - \mathbf{b} \rangle + 2t\langle \mathbf{v}, A\mathbf{v} \rangle,$$

the minimum occurs when

$$\hat{t} = -\frac{\langle \mathbf{v}, A\mathbf{x} - \mathbf{b} \rangle}{\langle \mathbf{v}, A\mathbf{v} \rangle} = \frac{\langle \mathbf{v}, \mathbf{b} - A\mathbf{x} \rangle}{\langle \mathbf{v}, A\mathbf{v} \rangle},$$

and, from Equation (7.27),

$$\begin{aligned} h(\hat{t}) &= g(\mathbf{x}) - 2\frac{\langle \mathbf{v}, \mathbf{b} - A\mathbf{x} \rangle}{\langle \mathbf{v}, A\mathbf{v} \rangle} \langle \mathbf{v}, \mathbf{b} - A\mathbf{x} \rangle + \left(\frac{\langle \mathbf{v}, \mathbf{b} - A\mathbf{x} \rangle}{\langle \mathbf{v}, A\mathbf{v} \rangle} \right)^2 \langle \mathbf{v}, A\mathbf{v} \rangle \\ &= g(\mathbf{x}) - \frac{\langle \mathbf{v}, \mathbf{b} - A\mathbf{x} \rangle^2}{\langle \mathbf{v}, A\mathbf{v} \rangle}. \end{aligned}$$

So, for any vector $\mathbf{v} \neq \mathbf{0}$, we have $g(\mathbf{x} + \hat{t}\mathbf{v}) < g(\mathbf{x})$ unless $\langle \mathbf{v}, \mathbf{b} - A\mathbf{x} \rangle = 0$, in which case $g(\mathbf{x}) = g(\mathbf{x} + \hat{t}\mathbf{v})$. This is the basic result we need to prove Theorem 7.31.

Suppose \mathbf{x}^* satisfies $A\mathbf{x}^* = \mathbf{b}$. Then $\langle \mathbf{v}, \mathbf{b} - A\mathbf{x}^* \rangle = 0$ for any vector \mathbf{v} , and $g(\mathbf{x})$ cannot be made any smaller than $g(\mathbf{x}^*)$. Thus, \mathbf{x}^* minimizes g .

On the other hand, suppose that \mathbf{x}^* is a vector that minimizes g . Then for any vector \mathbf{v} , we have $g(\mathbf{x}^* + \hat{t}\mathbf{v}) \geq g(\mathbf{x}^*)$. Thus, $\langle \mathbf{v}, \mathbf{b} - A\mathbf{x}^* \rangle = 0$. This implies that $\mathbf{b} - A\mathbf{x}^* = \mathbf{0}$ and, consequently, that $A\mathbf{x}^* = \mathbf{b}$. ■ ■ ■

To begin the conjugate gradient method, we choose \mathbf{x} , an approximate solution to $A\mathbf{x} = \mathbf{b}$, and $\mathbf{v} \neq \mathbf{0}$, which gives a *search direction* in which to move away from \mathbf{x} to improve the approximation. Let $\mathbf{r} = \mathbf{b} - A\mathbf{x}$ be the residual vector associated with \mathbf{x} and

$$t = \frac{\langle \mathbf{v}, \mathbf{b} - A\mathbf{x} \rangle}{\langle \mathbf{v}, A\mathbf{v} \rangle} = \frac{\langle \mathbf{v}, \mathbf{r} \rangle}{\langle \mathbf{v}, A\mathbf{v} \rangle}.$$

If $\mathbf{r} \neq \mathbf{0}$ and if \mathbf{v} and \mathbf{r} are not orthogonal, then $\mathbf{x} + t\mathbf{v}$ gives a smaller value for g than $g(\mathbf{x})$ and is presumably closer to \mathbf{x}^* than is \mathbf{x} . This suggests the following method.

Let $\mathbf{x}^{(0)}$ be an initial approximation to \mathbf{x}^* , and let $\mathbf{v}^{(1)} \neq \mathbf{0}$ be an initial search direction. For $k = 1, 2, 3, \dots$, we compute

$$t_k = \frac{\langle \mathbf{v}^{(k)}, \mathbf{b} - A\mathbf{x}^{(k-1)} \rangle}{\langle \mathbf{v}^{(k)}, A\mathbf{v}^{(k)} \rangle},$$

$$\mathbf{x}^{(k)} = \mathbf{x}^{(k-1)} + t_k \mathbf{v}^{(k)}$$

and choose a new search direction $\mathbf{v}^{(k+1)}$. The object is to make this selection so that the sequence of approximations $\{\mathbf{x}^{(k)}\}$ converges rapidly to \mathbf{x}^* .

To choose the search directions, we view g as a function of the components of $\mathbf{x} = (x_1, x_2, \dots, x_n)^t$. Thus,

(7.27)

$$g(x_1, x_2, \dots, x_n) = \langle \mathbf{x}, A\mathbf{x} \rangle - 2\langle \mathbf{x}, \mathbf{b} \rangle = \sum_{i=1}^n \sum_{j=1}^n a_{ij} x_i x_j - 2 \sum_{i=1}^n x_i b_i.$$

Taking partial derivatives with respect to the component variables x_k gives

$$\frac{\partial g}{\partial x_k}(\mathbf{x}) = 2 \sum_{i=1}^n a_{ki} x_i - 2b_k.$$

Therefore, the gradient of g is

$$\nabla g(\mathbf{x}) = \left(\frac{\partial g}{\partial x_1}(\mathbf{x}), \frac{\partial g}{\partial x_2}(\mathbf{x}), \dots, \frac{\partial g}{\partial x_n}(\mathbf{x}) \right)^t = 2(A\mathbf{x} - \mathbf{b}) = -2\mathbf{r},$$

where the vector \mathbf{r} is the residual vector for \mathbf{x} .

From multivariable calculus, we know that the direction of greatest decrease in the value of $g(\mathbf{x})$ is the direction given by $-\nabla g(\mathbf{x})$; that is, in the direction of the residual \mathbf{r} . The method that chooses

$$\mathbf{v}^{(k+1)} = \mathbf{r}^{(k)} = \mathbf{b} - A\mathbf{x}^{(k)}$$

is called the *method of steepest descent*. Although we will see in Section 10.4 that this method has merit for nonlinear systems and optimization problems, it is not used for linear systems because of slow convergence.

An alternative approach uses a set of nonzero direction vectors $\{\mathbf{v}^{(1)}, \dots, \mathbf{v}^{(n)}\}$ that satisfy

$$\langle \mathbf{v}^{(i)}, A\mathbf{v}^{(j)} \rangle = 0, \quad \text{if } i \neq j.$$

This is called an ***A-orthogonality condition***, and the set of vectors $\{\mathbf{v}^{(1)}, \dots, \mathbf{v}^{(n)}\}$ is said to be ***A-orthogonal***. It is not difficult to show that a set of *A-orthogonal* vectors associated with the positive definite matrix A is linearly independent. (See Exercise 13(a).) This set of search directions gives

$$t_k = \frac{\langle \mathbf{v}^{(k)}, \mathbf{b} - A\mathbf{x}^{(k-1)} \rangle}{\langle \mathbf{v}^{(k)}, A\mathbf{v}^{(k)} \rangle} = \frac{\langle \mathbf{v}^{(k)}, \mathbf{r}^{(k-1)} \rangle}{\langle \mathbf{v}^{(k)}, A\mathbf{v}^{(k)} \rangle}$$

and $\mathbf{x}^{(k)} = \mathbf{x}^{(k-1)} + t_k \mathbf{v}^{(k)}$.

The following theorem shows that this choice of search directions gives convergence in at most n -steps, so as a direct method it produces the exact solution, assuming that the arithmetic is exact.

Theorem 7.32

Let $\{\mathbf{v}^{(1)}, \dots, \mathbf{v}^{(n)}\}$ be an A -orthogonal set of nonzero vectors associated with the positive definite matrix A , and let $\mathbf{x}^{(0)}$ be arbitrary. Define

$$t_k = \frac{\langle \mathbf{v}^{(k)}, \mathbf{b} - A\mathbf{x}^{(k-1)} \rangle}{\langle \mathbf{v}^{(k)}, A\mathbf{v}^{(k)} \rangle} \quad \text{and} \quad \mathbf{x}^{(k)} = \mathbf{x}^{(k-1)} + t_k \mathbf{v}^{(k)},$$

for $k = 1, 2, \dots, n$. Then, assuming exact arithmetic, $A\mathbf{x}^{(n)} = \mathbf{b}$. ■

Proof Since, for each $k = 1, 2, \dots, n$,

$$\mathbf{x}^{(k)} = \mathbf{x}^{(k-1)} + t_k \mathbf{v}^{(k)},$$

we have

$$\begin{aligned} A\mathbf{x}^{(n)} &= A\mathbf{x}^{(n-1)} + t_n A\mathbf{v}^{(n)} \\ &= (A\mathbf{x}^{(n-2)} + t_{n-1} A\mathbf{v}^{(n-1)}) + t_n A\mathbf{v}^{(n)} \\ &\quad \vdots \\ &= A\mathbf{x}^{(0)} + t_1 A\mathbf{v}^{(1)} + t_2 A\mathbf{v}^{(2)} + \dots + t_n A\mathbf{v}^{(n)}, \end{aligned}$$

and subtracting \mathbf{b} from this result yields

$$A\mathbf{x}^{(n)} - \mathbf{b} = A\mathbf{x}^{(0)} - \mathbf{b} + t_1 A\mathbf{v}^{(1)} + t_2 A\mathbf{v}^{(2)} + \dots + t_n A\mathbf{v}^{(n)}.$$

We now take the inner product of both sides with the vector $\mathbf{v}^{(k)}$ and use the properties of inner products and the fact that A is symmetric to obtain

$$\begin{aligned} \langle A\mathbf{x}^{(n)} - \mathbf{b}, \mathbf{v}^{(k)} \rangle &= \langle A\mathbf{x}^{(0)} - \mathbf{b}, \mathbf{v}^{(k)} \rangle + t_1 \langle A\mathbf{v}^{(1)}, \mathbf{v}^{(k)} \rangle + \dots + t_n \langle A\mathbf{v}^{(n)}, \mathbf{v}^{(k)} \rangle \\ &= \langle A\mathbf{x}^{(0)} - \mathbf{b}, \mathbf{v}^{(k)} \rangle + t_1 \langle \mathbf{v}^{(1)}, A\mathbf{v}^{(k)} \rangle + \dots + t_n \langle \mathbf{v}^{(n)}, A\mathbf{v}^{(k)} \rangle. \end{aligned}$$

The A -orthogonality property gives, for each k ,

$$\langle A\mathbf{x}^{(n)} - \mathbf{b}, \mathbf{v}^{(k)} \rangle = \langle A\mathbf{x}^{(0)} - \mathbf{b}, \mathbf{v}^{(k)} \rangle + t_k \langle \mathbf{v}^{(k)}, A\mathbf{v}^{(k)} \rangle. \quad (7.28)$$

However,

$$t_k = \frac{\langle \mathbf{v}^{(k)}, \mathbf{b} - A\mathbf{x}^{(k-1)} \rangle}{\langle \mathbf{v}^{(k)}, A\mathbf{v}^{(k)} \rangle},$$

so

$$\begin{aligned} t_k \langle \mathbf{v}^{(k)}, A\mathbf{v}^{(k)} \rangle &= \langle \mathbf{v}^{(k)}, \mathbf{b} - A\mathbf{x}^{(k-1)} \rangle \\ &= \langle \mathbf{v}^{(k)}, \mathbf{b} - A\mathbf{x}^{(0)} + A\mathbf{x}^{(0)} - A\mathbf{x}^{(1)} + \dots - A\mathbf{x}^{(k-2)} + A\mathbf{x}^{(k-2)} - A\mathbf{x}^{(k-1)} \rangle \\ &= \langle \mathbf{v}^{(k)}, \mathbf{b} - A\mathbf{x}^{(0)} \rangle + \langle \mathbf{v}^{(k)}, A\mathbf{x}^{(0)} - A\mathbf{x}^{(1)} \rangle + \dots + \langle \mathbf{v}^{(k)}, A\mathbf{x}^{(k-2)} - A\mathbf{x}^{(k-1)} \rangle. \end{aligned}$$

But for any i ,

$$\mathbf{x}^{(i)} = \mathbf{x}^{(i-1)} + t_i \mathbf{v}^{(i)} \quad \text{and} \quad A\mathbf{x}^{(i)} = A\mathbf{x}^{(i-1)} + t_i A\mathbf{v}^{(i)},$$

so

$$A\mathbf{x}^{(i-1)} - A\mathbf{x}^{(i)} = -t_i A\mathbf{v}^{(i)}.$$

Thus,

$$t_k \langle \mathbf{v}^{(k)}, A\mathbf{v}^{(k)} \rangle = \langle \mathbf{v}^{(k)}, \mathbf{b} - A\mathbf{x}^{(0)} \rangle - t_1 \langle \mathbf{v}^{(k)}, A\mathbf{v}^{(1)} \rangle - \dots - t_{k-1} \langle \mathbf{v}^{(k)}, A\mathbf{v}^{(k-1)} \rangle.$$

Because of the A -orthogonality, $\langle \mathbf{v}^{(k)}, A\mathbf{v}^{(i)} \rangle = 0$, for $i \neq k$, so

$$\langle \mathbf{v}^{(k)}, A\mathbf{v}^{(k)} \rangle t_k = \langle \mathbf{v}^{(k)}, \mathbf{b} - A\mathbf{x}^{(0)} \rangle.$$

From Eq.(7.28),

$$\begin{aligned} \langle A\mathbf{x}^{(n)} - \mathbf{b}, \mathbf{v}^{(k)} \rangle &= \langle A\mathbf{x}^{(0)} - \mathbf{b}, \mathbf{v}^{(k)} \rangle + \langle \mathbf{v}^{(k)}, \mathbf{b} - A\mathbf{x}^{(0)} \rangle \\ &= \langle A\mathbf{x}^{(0)} - \mathbf{b}, \mathbf{v}^{(k)} \rangle + \langle \mathbf{b} - A\mathbf{x}^{(0)}, \mathbf{v}^{(k)} \rangle \\ &= \langle A\mathbf{x}^{(0)} - \mathbf{b}, \mathbf{v}^{(k)} \rangle - \langle A\mathbf{x}^{(0)} - \mathbf{b}, \mathbf{v}^{(k)} \rangle \\ &= 0. \end{aligned}$$

The vector $A\mathbf{x}^{(n)} - \mathbf{b}$ is orthogonal to the A -orthogonal set of vectors $\{\mathbf{v}^{(1)}, \dots, \mathbf{v}^{(n)}\}$. From this, it follows (see Exercise 13(b)) that $A\mathbf{x}^{(n)} - \mathbf{b} = \mathbf{0}$. ■ ■ ■

EXAMPLE 1 Consider the positive definite matrix

$$A = \begin{bmatrix} 4 & 3 & 0 \\ 3 & 4 & -1 \\ 0 & -1 & 4 \end{bmatrix}.$$

Let $\mathbf{v}^{(1)} = (1, 0, 0)^t$, $\mathbf{v}^{(2)} = (-3/4, 1, 0)^t$, and $\mathbf{v}^{(3)} = (-3/7, 4/7, 1)^t$. By direct calculation,

$$\langle \mathbf{v}^{(1)}, A\mathbf{v}^{(2)} \rangle = \mathbf{v}^{(1)t} A\mathbf{v}^{(2)} = (1, 0, 0) \begin{bmatrix} 4 & 3 & 0 \\ 3 & 4 & -1 \\ 0 & -1 & 4 \end{bmatrix} \begin{bmatrix} -3/4 \\ 1 \\ 0 \end{bmatrix} = 0,$$

$$\langle \mathbf{v}^{(1)}, A\mathbf{v}^{(3)} \rangle = (1, 0, 0) \begin{bmatrix} 4 & 3 & 0 \\ 3 & 4 & -1 \\ 0 & -1 & 4 \end{bmatrix} \begin{bmatrix} -3/7 \\ 4/7 \\ 1 \end{bmatrix} = 0,$$

and

$$\langle \mathbf{v}^{(2)}, A\mathbf{v}^{(3)} \rangle = \left(-\frac{3}{4}, 1, 0\right) \begin{bmatrix} 4 & 3 & 0 \\ 3 & 4 & -1 \\ 0 & -1 & 4 \end{bmatrix} \begin{bmatrix} -3/7 \\ 4/7 \\ 1 \end{bmatrix} = 0.$$

Thus, $\{\mathbf{v}^{(1)}, \mathbf{v}^{(2)}, \mathbf{v}^{(3)}\}$ is an A -orthogonal set.

The linear system

$$\begin{bmatrix} 4 & 3 & 0 \\ 3 & 4 & -1 \\ 0 & -1 & 4 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 24 \\ 30 \\ -24 \end{bmatrix},$$

has the exact solution $\mathbf{x}^* = (3, 4, -5)^t$. To approximate this solution, let $\mathbf{x}^{(0)} = (0, 0, 0)^t$. Since $\mathbf{b} = (24, 30, -24)^t$, we have

$$\mathbf{r}^{(0)} = \mathbf{b} - A\mathbf{x}^{(0)} = \mathbf{b} = (24, 30, -24)^t,$$

so

$$\langle \mathbf{v}^{(1)}, \mathbf{r}^{(0)} \rangle = \mathbf{v}^{(1)t} \mathbf{r}^{(0)} = 24, \quad \langle \mathbf{v}^{(1)}, A\mathbf{v}^{(1)} \rangle = 4, \quad \text{and} \quad t_0 = \frac{24}{4} = 6.$$

Thus,

$$\mathbf{x}^{(1)} = \mathbf{x}^{(0)} + t_0 \mathbf{v}^{(1)} = (0, 0, 0)^t + 6(1, 0, 0)^t = (6, 0, 0)^t.$$

Continuing, we have

$$\mathbf{r}^{(1)} = \mathbf{b} - A\mathbf{x}^{(1)} = (0, 12, -24)^t; \quad t_1 = \frac{\langle \mathbf{v}^{(2)}, \mathbf{r}^{(1)} \rangle}{\langle \mathbf{v}^{(2)}, A\mathbf{v}^{(2)} \rangle} = \frac{12}{7/4} = \frac{48}{7};$$

$$\mathbf{x}^{(2)} = \mathbf{x}^{(1)} + t_1 \mathbf{v}^{(2)} = (6, 0, 0)^t + \frac{48}{7} \left(-\frac{3}{4}, 1, 0 \right)^t = \left(\frac{6}{7}, \frac{48}{7}, 0 \right)^t;$$

$$\mathbf{r}^{(2)} = \mathbf{b} - A\mathbf{x}^{(2)} = \left(0, 0, -\frac{120}{7} \right)^t; \quad t_2 = \frac{\langle \mathbf{v}^{(3)}, \mathbf{r}^{(2)} \rangle}{\langle \mathbf{v}^{(3)}, A\mathbf{v}^{(3)} \rangle} = \frac{-120/7}{24/7} = -5;$$

and

$$\mathbf{x}^{(3)} = \mathbf{x}^{(2)} + t_2 \mathbf{v}^{(3)} = \left(\frac{6}{7}, \frac{48}{7}, 0 \right)^t + (-5) \left(-\frac{3}{7}, \frac{4}{7}, 1 \right)^t = (3, 4, -5)^t.$$

Since we applied the technique $n = 3$ times, this is the actual solution. ■

Before discussing how to determine the A -orthogonal set, we will continue the development. The use of an A -orthogonal set $\{\mathbf{v}^{(1)}, \dots, \mathbf{v}^{(n)}\}$ of direction vectors gives what is called a *conjugate direction* method. The following theorem shows the orthogonality of the residual vectors $\mathbf{r}^{(k)}$ and the direction vectors $\mathbf{v}^{(j)}$. A proof of this result using mathematical induction is considered in Exercise 14.

Theorem 7.33 The residual vectors $\mathbf{r}^{(k)}$, where $k = 1, 2, \dots, n$, for a conjugate direction method, satisfy the equations

$$\langle \mathbf{r}^{(k)}, \mathbf{v}^{(j)} \rangle = 0, \quad \text{for each } j = 1, 2, \dots, k. \quad \blacksquare$$

The conjugate gradient method of Hestenes and Stiefel chooses the search directions $\{\mathbf{v}^{(k)}\}$ during the iterative process so that the residual vectors $\{\mathbf{r}^{(k)}\}$ are mutually orthogonal. To construct the direction vectors $\{\mathbf{v}^{(1)}, \mathbf{v}^{(2)}, \dots\}$ and the approximations $\{\mathbf{x}^{(1)}, \mathbf{x}^{(2)}, \dots\}$, we start with an initial approximation $\mathbf{x}^{(0)}$ and use the steepest descent direction $\mathbf{r}^{(0)} = \mathbf{b} - A\mathbf{x}^{(0)}$ as the first search direction $\mathbf{v}^{(1)}$.

Assume that the conjugate directions $\mathbf{v}^{(1)}, \dots, \mathbf{v}^{(k-1)}$ and the approximations $\mathbf{x}^{(1)}, \dots, \mathbf{x}^{(k-1)}$ have been computed with

$$\mathbf{x}^{(k-1)} = \mathbf{x}^{(k-2)} + t_{k-1}\mathbf{v}^{(k-1)},$$

where

$$\langle \mathbf{v}^{(i)}, A\mathbf{v}^{(j)} \rangle = 0 \quad \text{and} \quad \langle \mathbf{r}^{(i)}, \mathbf{r}^{(j)} \rangle = 0, \quad \text{for } i \neq j.$$

If $\mathbf{x}^{(k-1)}$ is the solution to $A\mathbf{x} = \mathbf{b}$, we are done. Otherwise, $\mathbf{r}^{(k-1)} = \mathbf{b} - A\mathbf{x}^{(k-1)} \neq \mathbf{0}$ and Theorem 7.33 implies that $\langle \mathbf{r}^{(k-1)}, \mathbf{v}^{(i)} \rangle = 0$, for $i = 1, 2, \dots, k-1$. We then use $\mathbf{r}^{(k-1)}$ to generate $\mathbf{v}^{(k)}$ by setting

$$\mathbf{v}^{(k)} = \mathbf{r}^{(k-1)} + s_{k-1}\mathbf{v}^{(k-1)}.$$

We want to choose s_{k-1} so that

$$\langle \mathbf{v}^{(k-1)}, A\mathbf{v}^{(k)} \rangle = 0.$$

Since

$$A\mathbf{v}^{(k)} = A\mathbf{r}^{(k-1)} + s_{k-1}A\mathbf{v}^{(k-1)}$$

and

$$\langle \mathbf{v}^{(k-1)}, A\mathbf{v}^{(k)} \rangle = \langle \mathbf{v}^{(k-1)}, A\mathbf{r}^{(k-1)} \rangle + s_{k-1}\langle \mathbf{v}^{(k-1)}, A\mathbf{v}^{(k-1)} \rangle,$$

we will have $\langle \mathbf{v}^{(k-1)}, A\mathbf{v}^{(k)} \rangle = 0$ when

$$s_{k-1} = -\frac{\langle \mathbf{v}^{(k-1)}, A\mathbf{r}^{(k-1)} \rangle}{\langle \mathbf{v}^{(k-1)}, A\mathbf{v}^{(k-1)} \rangle}.$$

It can also be shown that with this choice of s_{k-1} we have $\langle \mathbf{v}^{(k)}, A\mathbf{v}^{(i)} \rangle = 0$, for each $i = 1, 2, \dots, k-2$ (see [Lu, p. 245]). Thus, $\{\mathbf{v}^{(1)}, \dots, \mathbf{v}^{(k)}\}$ is an A -orthogonal set.

Having chosen $\mathbf{v}^{(k)}$, we compute

$$\begin{aligned} t_k &= \frac{\langle \mathbf{v}^{(k)}, \mathbf{r}^{(k-1)} \rangle}{\langle \mathbf{v}^{(k)}, A\mathbf{v}^{(k)} \rangle} = \frac{\langle \mathbf{r}^{(k-1)} + s_{k-1}\mathbf{v}^{(k-1)}, \mathbf{r}^{(k-1)} \rangle}{\langle \mathbf{v}^{(k)}, A\mathbf{v}^{(k)} \rangle} \\ &= \frac{\langle \mathbf{r}^{(k-1)}, \mathbf{r}^{(k-1)} \rangle}{\langle \mathbf{v}^{(k)}, A\mathbf{v}^{(k)} \rangle} + s_{k-1} \frac{\langle \mathbf{v}^{(k-1)}, \mathbf{r}^{(k-1)} \rangle}{\langle \mathbf{v}^{(k)}, A\mathbf{v}^{(k)} \rangle}. \end{aligned}$$

By Theorem 7.33, $\langle \mathbf{v}^{(k-1)}, \mathbf{r}^{(k-1)} \rangle = 0$, so

$$t_k = \frac{\langle \mathbf{r}^{(k-1)}, \mathbf{r}^{(k-1)} \rangle}{\langle \mathbf{v}^{(k)}, A\mathbf{v}^{(k)} \rangle}. \quad (7.29)$$

Thus,

$$\mathbf{x}^{(k)} = \mathbf{x}^{(k-1)} + t_k \mathbf{v}^{(k)}.$$

To compute $\mathbf{r}^{(k)}$, we multiply by A and subtract \mathbf{b} to obtain

$$A\mathbf{x}^{(k)} - \mathbf{b} = A\mathbf{x}^{(k-1)} - \mathbf{b} + t_k A\mathbf{v}^{(k)}$$

or

$$\mathbf{r}^{(k)} = \mathbf{r}^{(k-1)} - t_k A\mathbf{v}^{(k)}.$$

Thus,

$$\langle \mathbf{r}^{(k)}, \mathbf{r}^{(k)} \rangle = \langle \mathbf{r}^{(k-1)}, \mathbf{r}^{(k)} \rangle - t_k \langle A\mathbf{v}^{(k)}, \mathbf{r}^{(k)} \rangle = -t_k \langle \mathbf{r}^{(k)}, A\mathbf{v}^{(k)} \rangle.$$

Further, from Eq. (7.29),

$$t_k = \frac{\langle \mathbf{r}^{(k-1)}, \mathbf{r}^{(k-1)} \rangle}{\langle \mathbf{v}^{(k)}, A\mathbf{v}^{(k)} \rangle}$$

so

$$\begin{aligned} s_k &= -\frac{\langle \mathbf{v}^{(k)}, A\mathbf{r}^{(k)} \rangle}{\langle \mathbf{v}^{(k)}, A\mathbf{v}^{(k)} \rangle} = -\frac{\langle \mathbf{r}^{(k)}, A\mathbf{v}^{(k)} \rangle}{\langle \mathbf{v}^{(k)}, A\mathbf{v}^{(k)} \rangle} \\ &= \frac{(1/t_k) \langle \mathbf{r}^{(k)}, \mathbf{r}^{(k)} \rangle}{(1/t_k) \langle \mathbf{r}^{(k-1)}, \mathbf{r}^{(k-1)} \rangle} = \frac{\langle \mathbf{r}^{(k)}, \mathbf{r}^{(k)} \rangle}{\langle \mathbf{r}^{(k-1)}, \mathbf{r}^{(k-1)} \rangle}. \end{aligned}$$

In summary, we have the formulas:

$$\mathbf{r}^{(0)} = \mathbf{b} - A\mathbf{x}^{(0)}; \quad \mathbf{v}^{(1)} = \mathbf{r}^{(0)};$$

and, for $k = 1, 2, \dots, n$,

$$\begin{aligned} t_k &= \frac{\langle \mathbf{r}^{(k-1)}, \mathbf{r}^{(k-1)} \rangle}{\langle \mathbf{v}^{(k)}, A\mathbf{v}^{(k)} \rangle}, \\ \mathbf{x}^{(k)} &= \mathbf{x}^{(k-1)} + t_k \mathbf{v}^{(k)}, \\ \mathbf{r}^{(k)} &= \mathbf{r}^{(k-1)} - t_k A\mathbf{v}^{(k)}, \\ s_k &= \frac{\langle \mathbf{r}^{(k)}, \mathbf{r}^{(k)} \rangle}{\langle \mathbf{r}^{(k-1)}, \mathbf{r}^{(k-1)} \rangle}, \\ \mathbf{v}^{(k+1)} &= \mathbf{r}^{(k)} + s_k \mathbf{v}^{(k)}. \end{aligned} \tag{7.30}$$

Rather than presenting an algorithm for the conjugate gradient method using these formulas, we extend the method to include *preconditioning*. If the matrix A is ill-conditioned, the conjugate gradient method is highly susceptible to rounding errors. So, although the exact answer should be obtained in n steps, this is not usually the case. As a direct method the conjugate gradient method is not as good as Gaussian elimination with pivoting. The main use of the conjugate gradient method is as an iterative method applied to a better-

conditioned system. In this case an acceptable approximate solution is often obtained in about \sqrt{n} steps.

To apply the method to a better-conditioned system, we want to select a nonsingular conditioning matrix C so that

$$\tilde{A} = C^{-1}A(C^{-1})^t$$

is better conditioned. To simplify the notation, we will use the matrix C^{-t} to refer to $(C^{-1})^t$.

Consider the linear system

$$\tilde{A}\tilde{\mathbf{x}} = \tilde{\mathbf{b}},$$

where $\tilde{\mathbf{x}} = C^t\mathbf{x}$ and $\tilde{\mathbf{b}} = C^{-1}\mathbf{b}$. Then

$$\tilde{A}\tilde{\mathbf{x}} = (C^{-1}AC^{-t})(C^t\mathbf{x}) = C^{-1}A\mathbf{x}.$$

Thus, we could solve $\tilde{A}\tilde{\mathbf{x}} = \tilde{\mathbf{b}}$ for $\tilde{\mathbf{x}}$ and then obtain \mathbf{x} by multiplying by C^{-t} . However, instead of rewriting equations (7.30) using $\tilde{\mathbf{r}}^{(k)}$, $\tilde{\mathbf{v}}^{(k)}$, \tilde{t}_k , $\tilde{\mathbf{x}}^{(k)}$, and \tilde{s}_k , we incorporate the preconditioning implicitly.

Since

$$\tilde{\mathbf{x}}^{(k)} = C^t\mathbf{x}^{(k)},$$

we have

$$\tilde{\mathbf{r}}^{(k)} = \tilde{\mathbf{b}} - \tilde{A}\tilde{\mathbf{x}}^{(k)} = C^{-1}\mathbf{b} - (C^{-1}AC^{-t})C^t\mathbf{x}^{(k)} = C^{-1}(\mathbf{b} - A\mathbf{x}^{(k)}) = C^{-1}\mathbf{r}^{(k)}.$$

Let $\tilde{\mathbf{v}}^{(k)} = C^t\mathbf{v}^{(k)}$ and $\tilde{\mathbf{w}}^{(k)} = C^{-1}\mathbf{r}^{(k)}$. Then

$$\tilde{s}_k = \frac{\langle \tilde{\mathbf{r}}^{(k)}, \tilde{\mathbf{r}}^{(k)} \rangle}{\langle \tilde{\mathbf{r}}^{(k-1)}, \tilde{\mathbf{r}}^{(k-1)} \rangle} = \frac{\langle C^{-1}\mathbf{r}^{(k)}, C^{-1}\mathbf{r}^{(k)} \rangle}{\langle C^{-1}\mathbf{r}^{(k-1)}, C^{-1}\mathbf{r}^{(k-1)} \rangle},$$

so

$$\tilde{s}_k = \frac{\langle \mathbf{w}^{(k)}, \mathbf{w}^{(k)} \rangle}{\langle \mathbf{w}^{(k-1)}, \mathbf{w}^{(k-1)} \rangle}. \quad (7.31) \quad (4)$$

Thus,

$$\tilde{t}_k = \frac{\langle \tilde{\mathbf{r}}^{(k-1)}, \tilde{\mathbf{r}}^{(k-1)} \rangle}{\langle \tilde{\mathbf{v}}^{(k)}, \tilde{A}\tilde{\mathbf{v}}^{(k)} \rangle} = \frac{\langle C^{-1}\mathbf{r}^{(k-1)}, C^{-1}\mathbf{r}^{(k-1)} \rangle}{\langle C^t\mathbf{v}^{(k)}, C^{-1}AC^{-t}C^t\mathbf{v}^{(k)} \rangle} = \frac{\langle \mathbf{w}^{(k-1)}, \mathbf{w}^{(k-1)} \rangle}{\langle C^t\mathbf{v}^{(k)}, C^{-1}A\mathbf{v}^{(k)} \rangle}$$

and

$$\tilde{t}_k = \frac{\langle \mathbf{w}^{(k-1)}, \mathbf{w}^{(k-1)} \rangle}{\langle \mathbf{v}^{(k)}, A\mathbf{v}^{(k)} \rangle}. \quad (7.32) \quad (1)$$

Further,

$$\tilde{\mathbf{x}}^{(k)} = \tilde{\mathbf{x}}^{(k-1)} + \tilde{t}_k\tilde{\mathbf{v}}^{(k)}, \quad \text{so} \quad C^t\mathbf{x}^{(k)} = C^t\mathbf{x}^{(k-1)} + \tilde{t}_kC^t\mathbf{v}^{(k)}$$

and

$$\mathbf{x}^{(k)} = \mathbf{x}^{(k-1)} + \tilde{t}_k\mathbf{v}^{(k)}. \quad (7.33) \quad (2)$$

(7.30)
 ese for-
 itioned,
 ough the
 method
 ng. The
 better-

Continuing,

$$\tilde{\mathbf{r}}^{(k)} = \tilde{\mathbf{r}}^{(k-1)} - \tilde{t}_k \tilde{A} \tilde{\mathbf{v}}^{(k)},$$

so

$$C^{-1} \mathbf{r}^{(k)} = C^{-1} \mathbf{r}^{(k-1)} - \tilde{t}_k C^{-1} A C^{-t} \tilde{\mathbf{v}}^{(k)}, \quad \mathbf{r}^{(k)} = \mathbf{r}^{(k-1)} - \tilde{t}_k A C^{-t} C^t \mathbf{v}^{(k)},$$

and

$$\mathbf{r}^{(k)} = \mathbf{r}^{(k-1)} - \tilde{t}_k A \mathbf{v}^{(k)}. \quad (3) \quad (7.34)$$

Finally,

$$\tilde{\mathbf{v}}^{(k+1)} = \tilde{\mathbf{r}}^{(k)} + \tilde{s}_k \tilde{\mathbf{v}}^{(k)} \quad \text{and} \quad C^t \mathbf{v}^{(k+1)} = C^{-1} \mathbf{r}^{(k)} + \tilde{s}_k C^t \mathbf{v}^{(k)},$$

so

$$\mathbf{v}^{(k+1)} = C^{-t} C^{-1} \mathbf{r}^{(k)} + \tilde{s}_k \mathbf{v}^{(k)} = C^{-t} \mathbf{w}^{(k)} + \tilde{s}_k \mathbf{v}^{(k)}. \quad (5) \quad (7.35)$$

The preconditioned conjugate gradient method is based on using equations (7.31)–(35) in the order (7.32), (7.33), (7.34), (7.31), (7.35). Algorithm 7.5 implements this procedure.

ALGORITHM

7.5

Preconditioned Conjugate Gradient Method

To solve $A\mathbf{x} = \mathbf{b}$ given the preconditioning matrix C^{-1} and the initial approximation $\mathbf{x}^{(0)}$:

INPUT the number of equations and unknowns n ; the entries a_{ij} , $1 \leq i, j \leq n$ of the matrix A ; the entries b_j , $1 \leq j \leq n$ of the vector \mathbf{b} ; the entries γ_{ij} , $1 \leq i, j \leq n$ of the preconditioning matrix C^{-1} , the entries x_i , $1 \leq i \leq n$ of the initial approximation $\mathbf{x} = \mathbf{x}^{(0)}$, the maximum number of iterations N ; tolerance TOL .

OUTPUT the approximate solution x_1, \dots, x_n and the residual r_1, \dots, r_n or a message that the number of iterations was exceeded.

Step 1 Set $\mathbf{r} = \mathbf{b} - A\mathbf{x}$; (Compute $\mathbf{r}^{(0)}$.)
 $\mathbf{w} = C^{-1} \mathbf{r}$; (Note: $\mathbf{w} = \mathbf{w}^{(0)}$)
 $\mathbf{v} = C^{-t} \mathbf{w}$; (Note: $\mathbf{v} = \mathbf{v}^{(1)}$)
 $\alpha = \sum_{j=1}^n w_j^2$.

Step 2 Set $k = 1$.

Step 3 While ($k \leq N$) do Steps 4–7.

Step 4 If $\|\mathbf{v}\| < TOL$, then
 OUTPUT ('Solution vector'; x_1, \dots, x_n);
 OUTPUT ('with residual'; r_1, \dots, r_n);
 (The procedure was successful.)
 STOP

Step 5 Set $\mathbf{u} = A\mathbf{v}$; (Note: $\mathbf{u} = A\mathbf{v}^{(k)}$)
 $t = \frac{\alpha}{\sum_{j=1}^n v_j u_j}$; (Note: $t = t_k$)

$$\begin{aligned} \mathbf{x} &= \mathbf{x} + t\mathbf{v}; \text{ (Note: } \mathbf{x} = \mathbf{x}^{(k)}) \\ \mathbf{r} &= \mathbf{r} - t\mathbf{u}; \text{ (Note: } \mathbf{r} = \mathbf{r}^{(k)}) \\ \mathbf{w} &= C^{-1}\mathbf{r}; \text{ (Note: } \mathbf{w} = \mathbf{w}^{(k)}) \\ \beta &= \sum_{j=1}^n w_j^2. \text{ (Note: } \beta = \langle \mathbf{w}^{(k)}, \mathbf{w}^{(k)} \rangle) \end{aligned}$$

Step 6 If $|\beta| < TOL$ then
 if $\|\mathbf{r}\| < TOL$ then
 OUTPUT('Solution vector'; x_1, \dots, x_n);
 OUTPUT('with residual'; r_1, \dots, r_n);
 (The procedure was successful.)
 STOP

Step 7 Set $s = \beta/\alpha$; ($s = s_k$)
 $\mathbf{v} = C^{-t}\mathbf{w} + s\mathbf{v}$; (Note: $\mathbf{v} = \mathbf{v}^{(k+1)}$)
 $\alpha = \beta$; (Update α .)
 $k = k + 1$.

Step 8 If ($k > n$) then
 OUTPUT ('The maximum number of iterations was exceeded.');

(The procedure was unsuccessful.)
 STOP.

The next example illustrates the calculations in an easy problem.

EXAMPLE 2 The linear system $A\mathbf{x} = \mathbf{b}$ given by

$$\begin{aligned} 4x_1 + 3x_2 &= 24, \\ 3x_1 + 4x_2 - x_3 &= 30, \\ -x_2 + 4x_3 &= -24 \end{aligned}$$

has solution $(3, 4, -5)^t$ and was considered in Example 3 of Section 7.3. In that example, both the Gauss-Seidel method and SOR method were used. We will use the conjugate gradient method with no preconditioning, so $C = C^{-1} = I$. Let $\mathbf{x}^{(0)} = (0, 0, 0)^t$. Then

$$\begin{aligned} \mathbf{r}^{(0)} &= \mathbf{b} - A\mathbf{x}^{(0)} = \mathbf{b} = (24, 30, -24)^t; \\ \mathbf{w} &= C^{-1}\mathbf{r}^{(0)} = (24, 30, -24)^t; \\ \mathbf{v}^{(1)} &= C^{-t}\mathbf{w} = (24, 30, -24)^t; \\ \alpha &= \langle \mathbf{w}, \mathbf{w} \rangle = 2052. \end{aligned}$$

We start the first iteration with $k = 1$. Then

$$\begin{aligned} \mathbf{u} &= A\mathbf{v}^{(1)} = (186.0, 216.0, -126.0)^t; \\ t_1 &= \frac{\alpha}{\langle \mathbf{v}^{(1)}, \mathbf{u} \rangle} = 0.1469072165; \\ \mathbf{x}^{(1)} &= \mathbf{x}^{(0)} + t_1\mathbf{v}^{(1)} = (3.525773196, 4.407216495, -3.525773196)^t; \end{aligned}$$

$$\begin{aligned} \mathbf{r}^{(1)} &= \mathbf{r}^{(0)} - t_1 \mathbf{u} = (-3.32474227, -1.73195876, -5.48969072)^t; \\ \mathbf{w} &= C^{-1} \mathbf{r}^{(1)} = \mathbf{r}^{(1)}; \\ \beta &= \langle \mathbf{w}, \mathbf{w} \rangle = 44.19029651; \\ s_1 &= \frac{\beta}{\alpha} = 0.02153523222; \\ \mathbf{v}^{(2)} &= C^{-t} \mathbf{w} + s_1 \mathbf{v}^{(1)} = (-2.807896697, -1.085901793, -6.006536293)^t. \end{aligned}$$

Set

$$\alpha = \beta = 44.19029651.$$

We are now ready to begin the second iteration. We have

$$\begin{aligned} \mathbf{u} &= A\mathbf{v}^{(2)} = (-14.48929217, -6.760760967, -22.94024338)^t; \\ t_2 &= 0.2378157558; \\ \mathbf{x}^{(2)} &= (2.858011121, 4.148971939, -4.954222164)^t; \\ \mathbf{r}^{(2)} &= (0.121039698, -0.124143281, -0.034139402)^t; \\ \mathbf{w} &= C^{-1} \mathbf{r}^{(2)} = \mathbf{r}^{(2)}; \\ \beta &= 0.03122766148; \\ s_2 &= 0.0007066633163; \\ \mathbf{v}^{(3)} &= (0.1190554504, -0.1249106480, -0.03838400086)^t. \end{aligned}$$

Set

$$\alpha = \beta = 0.03122766148.$$

Finally, the third iteration gives

$$\begin{aligned} \mathbf{u} &= A\mathbf{v}^{(3)} = (0.1014898976, -0.1040922099, -0.0286253554)^t; \\ t_3 &= 1.192628008; \\ \mathbf{x}^{(3)} &= (2.999999998, 4.000000002, -4.999999998)^t; \\ \mathbf{r}^{(3)} &= (0.36 \times 10^{-8}, 0.39 \times 10^{-8}, -0.141 \times 10^{-8})^t. \end{aligned}$$

Since $\mathbf{x}^{(3)}$ is nearly the exact solution, rounding error did not significantly effect the result. In Example 3 of Section 7.3, the Gauss-Seidel method required 34 iterations, and the SOR method, with $\omega = 1.25$, required 14 iterations for an accuracy of 10^{-7} . It should be noted, however, that in this example, we are really comparing a direct method to iterative methods. ■

The next example illustrates the effect of preconditioning on a poorly conditioned matrix. In this example and subsequently, we use $D^{-1/2}$ to represent the diagonal matrix

whose e
matrix A

EXAMPLE 3 The line

has the s

$$\mathbf{x}^* = (7.$$

The mat
ber K_∞
the Jaco
jugate gr
the diag
of the di
7.5. The
with the

Table 7.5

Method	o
Jacobi	
Gauss-Seidel	
SOR ($\omega = 1.25$)	
Conjugate Gradient	
Conjugate Gradient (Preconditioned)	

The
linear sy
be solve
equation
jugate g
required
Cholesk
method
 $C^{-t} C^{-1}$
gate gra

whose entries are the reciprocals of the square roots of the diagonal entries of the coefficient matrix A .

EXAMPLE 3 The linear system $Ax = b$ with

$$A = \begin{bmatrix} 0.2 & 0.1 & 1 & 1 & 0 \\ 0.1 & 4 & -1 & 1 & -1 \\ 1 & -1 & 60 & 0 & -2 \\ 1 & 1 & 0 & 8 & 4 \\ 0 & -1 & -2 & 4 & 700 \end{bmatrix} \quad \text{and} \quad b = \begin{bmatrix} 1 \\ 2 \\ 3 \\ 4 \\ 5 \end{bmatrix}$$

has the solution

$$x^* = (7.859713071, 0.4229264082, -0.07359223906, -0.5406430164, 0.01062616286)^t.$$

The matrix A is symmetric and positive definite but is ill-conditioned with condition number $K_\infty(A) = 13961.71$. We will use tolerance 0.01 and compare the results obtained from the Jacobi, Gauss-Seidel, and SOR (with $\omega = 1.25$) iterative methods and from the conjugate gradient method with $C^{-1} = I$. Then we precondition by choosing C^{-1} as $D^{-1/2}$, the diagonal matrix whose diagonal entries are the reciprocal of the positive square roots of the diagonal entries of the positive definite matrix A . The results are presented in Table 7.5. The preconditioned conjugate gradient method gives the most accurate approximation with the smallest number of iterations. ■

Table 7.5

Method	Number of Iterations	$x^{(k)}$	$\ x^* - x^{(k)}\ _\infty$
Jacobi	49	(7.86277141, 0.42320802, -0.07348669, -0.53975964, 0.01062847) ^t	0.00305834
Gauss-Seidel	15	(7.83525748, 0.42257868, -0.07319124, -0.53753055, 0.01060903) ^t	0.02445559
SOR ($\omega = 1.25$)	7	(7.85152706, 0.42277371, -0.07348303, -0.53978369, 0.01062286) ^t	0.00818607
Conjugate Gradient	5	(7.85341523, 0.42298677, -0.07347963, -0.53987920, 0.008628916) ^t	0.00629785
Conjugate Gradient (Preconditioned)	4	(7.85968827, 0.42288329, -0.07359878, -0.54063200, 0.01064344) ^t	0.00009312

The preconditioned conjugate gradient method is often used in the solution of large linear systems in which the matrix is sparse and positive definite. These systems must be solved to approximate solutions to boundary-value problems in ordinary-differential equations (Sections 11.3, 11.4, 11.5). The larger the system, the more impressive the conjugate gradient method becomes since it significantly reduces the number of iterations required. In these systems, the preconditioning matrix C is approximately equal to L in the Choleski factorization LL' of A . Generally, small entries in A are ignored and Choleski's method is applied to obtain what is called an incomplete LL' factorization of A . Thus, $C^{-1}C^{-1} \approx A^{-1}$ and a good approximation is obtained. More information about the conjugate gradient method can be found in Kelley [Kelley].

$$C^{-1} = D^{-1/2}$$

).

y effect the
ons, and the
t should be
to iterative

conditioned
onal matrix