

% back substitution

%  $x(n) = b(n) / a(n,n);$

for  $i = n-1 : -1 : 1$

$s = 0;$

for  $j = i+1 : n$

$s = s + a(i,j) * x(j);$

end

$x(i) = (b(i) - s) / a(i,i);$

end

Note  $a_{11}x_1 + \underbrace{a_{12}x_2 + \dots + a_{1n}x_n}_s = b_1$

Operation count (reduction to upper  $\Delta$  form)

# divisions =  $1 + 2 + \dots + (n-1) = \frac{n(n-1)}{2}$

# multiplications =  $n(n-1)(2n-1) / 6$

$i = 1 : N \Rightarrow i = 1 : 1 : N$

$i = 1 : 2 : N$   
(step)

Pf # divisions =  $\sum_{k=1}^{n-1} (n-k) = \sum_{k=1}^{n-1} k = \frac{n(n-1)}{2}$

# mult =  $\sum_{k=1}^{n-1} (n-k)^2 = \sum_{k=1}^{n-1} k^2 = S$

$$\begin{aligned} (n-1)^3 &= (n-1)^3 - (n-2)^3 + (n-2)^3 + \dots - 2^3 + 2^3 - 1^3 + 1^3 = \\ &= \sum_{k=1}^{n-1} (k^3 - (k-1)^3) = \sum_{k=1}^{n-1} (\cancel{k^3} - \cancel{(k-1)^3}) = \\ &= 3 \underbrace{\sum_{k=1}^{n-1} k^2}_S - 3 \underbrace{\sum_{k=1}^{n-1} k}_{\frac{n(n-1)}{2}} + \underbrace{\sum_{k=1}^{n-1} 1}_{n-1} \end{aligned}$$

$$\Rightarrow (n-1)^3 = 3S - 3 \frac{n(n-1)}{2} + (n-1)$$

Solve for S.

$$S = \frac{n(n-1)(2n-1)}{6}$$





where

$$M_k = \begin{pmatrix} 0 & & & \\ \vdots & & & \\ 0 & & & \\ m_{k+1,k} & & & \\ \vdots & & & \\ m_{nk} & & & \end{pmatrix}$$

$$e_k = \begin{pmatrix} 0 & & & \\ \vdots & & & \\ 1 & & & \\ 0 & & & \\ \vdots & & & \\ 0 & & & \end{pmatrix} \leftarrow k^{\text{th}} \text{ row}$$

Ex  $3 \times 3$  case

$$E_1 = \begin{pmatrix} 1 & 0 & 0 \\ -m_{21} & 1 & 0 \\ -m_{31} & 0 & 1 \end{pmatrix} = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix} - \begin{pmatrix} 0 & 0 & 0 \\ m_{21} & 0 & 0 \\ m_{31} & 0 & 0 \end{pmatrix}$$

$$\begin{pmatrix} 0 & 0 & 0 \\ m_{21} & 0 & 0 \\ m_{31} & 0 & 0 \end{pmatrix} = \begin{pmatrix} 0 \\ m_{21} \\ m_{31} \end{pmatrix} (1 \ 0 \ 0)$$









1. Find  $L, U$  such that  $A = LU$   $O(n^3)$
2. Solve  $Ly = b$  for  $y$  } triangular systems  $O(n^2)$
3. Solve  $Ux = y$  for  $x$  }  $O(n^2)$

$$\left. \begin{array}{l} Ax = b \\ \underbrace{LU}x = b \\ y \end{array} \right\}$$

$$Ux = y$$

Then  $Ax = \underbrace{LU}x = Ly = b$

Advantage If you solve for  $L$  and  $U$ , then you can solve for many vectors  $b$  using steps 2 and 3.

Storage

Instead of storing  $L$  and  $U$ , you can overwrite  $A$  with reduced  $A$ , i.e.  $\tilde{U}$ , and entries of  $L$ :

